

## More on Square and Square Root of a Node on T<sub>3</sub> Tree

Xingbo WANG\*

Department of Mechatronics Engineering, Foshan University, PRC

Guangdong Engineering Center of Information Security for Intelligent Manufacturing System, PRC

State Key Laboratory of Mathematical Engineering and Advanced Computing, PRC

**\*Corresponding Author:** Xingbo WANG, State Key Laboratory of Mathematical Engineering and Advanced Computing, PRC, China.

**Abstract:** The article proves several inequalities derived from nodal multiplication on T<sub>3</sub> tree. The proved inequalities are helpful to estimate certain quantities related with the T<sub>3</sub> tree as well as examples of proving an inequality embedded with the floor functions.

**Keywords:** Square, Square root, Integer, Binary tree

### 1. INTRODUCTION

The square and the square root are undoubtedly very important operations for a number. As a new structure of odd integers, the T<sub>3</sub> tree, which was introduced in [1] and [2], is of course necessary to make clear these two operations. In fact, for a given node in the tree, the problem where its square and its square root locate is surely a fundamental problem. Paper [2] mentioned some general properties of the square of a node in T<sub>3</sub>, paper [4] was the first one to disclose the properties of the square root of a node. However, one can see that, there are still a lot of unknown properties. This paper shows a little more of the properties related with the square and the square root of a node.

### 2. PRELIMINARIES

This section lists for later sections the necessary preliminaries, which include definitions, notations and lemmas.

#### 2.1. Definitions and Notations

Let  $S$  be a set of finite positive integers with  $s_0$  and  $s_n$  being the smallest and the biggest nodes respectively; an integer  $x$  is said to be clamped in  $S$  if  $s_0 \leq x \leq s_n$ . Symbol  $x \in S$  indicates that  $x$  is clamped in  $S$ . Symbol  $\lfloor x \rfloor$  is the floor function, an integer function of real number  $x$  that satisfies inequality  $x - 1 < \lfloor x \rfloor \leq x$ , or equivalently  $\lfloor x \rfloor \leq x < \lfloor x \rfloor + 1$ .

In this whole paper, symbol  $T_3$  is the T<sub>3</sub> tree that was introduced in [1] and [2] and symbol  $N_{(k,j)}$  is by default the node at position  $j$  on level  $k$  of T<sub>3</sub>, where  $k \geq 0$  and  $0 \leq j \leq 2^k - 1$ . By using the asterisk wildcard \*, symbol  $N_{(k,*)}$  means a node lying on level  $k$ . An integer  $X$  is said to be clamped on level  $k$  of T<sub>3</sub> if  $2^{k+1} \leq X \leq 2^{k+2} - 1$  and symbol  $X \in k$  indicates  $X$  is clamped on level  $k$ . If a positive integer  $X$  is clamped on level  $k$  and there is a node  $Y$  of T<sub>3</sub> satisfying  $X = \lfloor \sqrt{Y} \rfloor$ , then  $X$  is said to be a floor square root of the node  $Y$  and  $Y$  is called a square source of  $X$ .

Remark 1. The concept that an integer is clamped on a level of T<sub>3</sub> was first put forward in paper [3]. In the paper [3], a positive integer  $X$  was said to be clamped on level  $k$  of T<sub>3</sub> if  $2^{k+1} + 1 \leq X \leq 2^{k+2} - 1$ . Since  $2^{k+1} - 1$  is the rightmost node on level  $k - 1$  and  $2^{k+1} + 1$  is the leftmost node on level  $k$ , there is an integer  $2^{k+1}$  between the two. In order to avoid leaving out the number  $2^{k+1}$ , this paper redefines it by  $2^{k+1} \leq X \leq 2^{k+2} - 1$ .

2.2. Lemmas

**Lemma 1 (See in [1]).**  $T_3$  Tree has the following fundamental properties.

**(P1).** Every node is an odd integer and every odd integer bigger than 1 must be on the  $T_3$  tree. Odd integer  $N$  with  $N > 1$  lies on level  $\lfloor \log_2 N \rfloor - 1$ .

**(P2).** On level  $k$  with  $k = 0, 1, \dots$ , there are  $2^k$  nodes starting by  $2^{k+1} + 1$  and ending by  $2^{k+2} - 1$ , namely,  $N_{(k,j)} \in [2^{k+1} + 1, 2^{k+2} - 1]$  with  $j = 0, 1, \dots, 2^k - 1$ .

**(P3).**  $N_{(k,j)}$  is calculated by

$$N_{(k,j)} = 2^{k+1} + 1 + 2j, j = 0, 1, \dots, 2^k - 1$$

**(P4)** Multiplication of arbitrary two nodes of  $T_3$ , say  $N_{(m,\alpha)}$  and  $N_{(n,\beta)}$ , is a third node of  $T_3$ . Let  $J = 2^m(1 + 2\beta) + 2^n(1 + 2\alpha) + 2\alpha\beta + \alpha + \beta$ ; the multiplication  $N_{(m,\alpha)} \times N_{(n,\beta)}$  is given by

$$N_{(m,\alpha)} \times N_{(n,\beta)} = 2^{m+n+2} + 1 + 2J$$

If  $J < 2^{m+n+1}$ , then  $N_{(m,\alpha)} \times N_{(n,\beta)} = N_{(m+n+1,J)}$  lies on level  $m+n+1$  of  $T_3$ ; whereas, if  $J \geq 2^{m+n+1}$ ,  $N_{(m,\alpha)} \times N_{(n,\beta)} = N_{(m+n+2,\chi)}$  with  $\chi = J - 2^{m+n+1}$  lies on level  $m+n+2$  of  $T_3$ .

**Lemma 2 (See in [4]).** For real numbers  $x$ ,  $y$  and positive integer  $i$ , it holds

**(P13)**  $x \leq y \Rightarrow \lfloor x \rfloor \leq \lfloor y \rfloor$ ;  $x < n \Rightarrow \lfloor x \rfloor < n$ , where  $n$  is an integer.

**(P31)**  $i - 1 \leq 2 \left\lfloor \frac{i}{2} \right\rfloor \leq i$ .

3. MAIN RESULTS

**Theorem 1.** Let  $k$  be a positive integer; then there are  $2k + 1$  consecutive integers  $n_1, n_2, \dots, n_{2k+1}$  that satisfy  $\lfloor \sqrt{n_i} \rfloor_{(i=1,2,\dots,2k+1)} = k$ .

**Proof.** Consider an arbitrary integer  $n$  such that  $\lfloor \sqrt{n} \rfloor = k$ ; then by definition of the floor function it holds  $k \leq \sqrt{n} < k + 1$ . That is

$$k^2 \leq n < k^2 + 2k + 1$$

Hence the  $2k + 1$  integers,  $n_1 = k^2 + 0$ ,  $n_2 = k^2 + 1, n_3 = k^2 + 2, \dots$ , and  $n_{2k+1} = k^2 + 2k$ , are the integers satisfying  $\lfloor \sqrt{n_i} \rfloor_{(i=1,2,\dots,2k+1)} = k$ .

**Proposition 1.** Let  $N_{(m,\alpha)}$  be a node of  $T_3$  with  $m > 0$ ; then when  $0 \leq \alpha \leq \left\lfloor \frac{\sqrt{2^{2m+3} + 1} - 1}{2} \right\rfloor - 2^m$ ,  $N_{(m,\alpha)}^2$  lies on level  $2m + 1$ ; otherwise it lies on level  $2m + 2$ . Particularly,  $N_{(m,0)}^2 = N_{(2m+1,2^{m+1})}$  and  $N_{(m,2^m-1)}^2 = N_{(2m+2,2^{2m+2}-2^{m+2})}$ .

**Proof.** Direct calculation shows

$$N_{(m,\alpha)}^2 = (2^{m+1} + 2\alpha + 1)^2 = 2^{2m+2} + 2(2\alpha^2 + 2^{m+2}\alpha + 2^{m+1} + 2\alpha) + 1$$

By Lemma 1 (**P4** & **P5**), it knows that  $N_{(m,\alpha)}^2$  lies on level  $2m + 1$  if and only if  $J = 2\alpha^2 + 2^{m+2}\alpha + 2^{m+1} + 2\alpha < 2^{2m+1}$ . Consequently,

$$\begin{aligned} \alpha^2 + 2^{m+1}\alpha + 2^m + \alpha &< 2^{2m} \\ \Rightarrow \alpha^2 + (2^{m+1} + 1)\alpha + 2^m - 2^{2m} &< 0 \\ \Rightarrow 0 \leq \alpha < \frac{-(2^{m+1} + 1) + \sqrt{(2^{m+1} + 1)^2 - 4 \times (2^m - 2^{2m})}}{2} \\ \Rightarrow 0 \leq \alpha < \frac{\sqrt{2^{2m+3} + 1} - (2^{m+1} + 1)}{2} \\ \Rightarrow 0 \leq \alpha < \frac{\sqrt{2^{2m+3} + 1} - 1}{2} - 2^m \\ \Rightarrow 0 \leq \alpha \leq \left\lfloor \frac{\sqrt{2^{2m+3} + 1} - 1}{2} \right\rfloor - 2^m \end{aligned}$$

which validates the first part of the proposition.

The second part is easily obtained by the following calculations.

$$\begin{aligned} N_{(m,0)}^2 &= (2^{m+1} + 2 \times 0 + 1)^2 = 2^{2m+2} + 2 \times 2^{m+1} + 1 = N_{(2m+1,2^{m+1})} \\ N_{(m,2^m-1)}^2 &= (2^{m+1} + 2 \times (2^m - 1) + 1)^2 = (2^{m+1} + 2^{m+1} - 1)^2 \\ &= 2^{2m+2} + 2^{2m+2} + 1 + 2^{2m+3} - 2^{m+2} - 2^{m+2} \\ &= 2^{2m+3} + 2^{2m+3} - 2^{m+3} + 1 \\ &= 2^{2m+3} + 2 \times (2^{2m+2} - 2^{m+2}) + 1 \\ &= N_{(2m+2,2^{2m+2}-2^{m+2})} \end{aligned}$$

**Remark 1.** The condition  $m > 0$  in Proposition 1 is proposed because it can get rid of the case  $N_{(0,0)}^2 = 3^2 = 9 = N_{(2,0)}$ , which is the unique example that violates  $N_{(m,0)}^2 = N_{(2m+1,2^{m+1})}$ .

**Proposition 2.** Let  $k$  be a positive integer,  $N_{(2k+1,*)}$  and  $N_{(2k+2,*)}$  be nodes on of  $T_3$ ; then  $\lfloor \sqrt{N_{(2k+1,*)}} \rfloor \leq \lfloor 2^{k+1}\sqrt{2} \rfloor \square k$  and  $\lfloor 2^{k+1}\sqrt{2} \rfloor \leq \lfloor \sqrt{N_{(2k+2,*)}} \rfloor \square k$ . On level  $2k$  there is not a node  $N_{(2k,*)}$  satisfying  $N_{(2k,*)} \square k$  and there is neither a node  $N_{(2k+3,*)}$  satisfying  $N_{(2k+3,*)} \square k$  on level  $2k + 3$ .

**Proof.**  $N_{(2k+1,*)}$  and  $N_{(2k+2,*)}$  being the nodes on levels  $2k + 1$  and  $2k + 2$  respectively yields

$$\begin{aligned} 2^{2k+2} < 2^{2k+2} + 1 \leq N_{(2k+1,*)} \leq 2^{2k+3} - 1 < 2^{2k+3} \\ \Rightarrow 2^{k+1} = \lfloor \sqrt{2^{2k+2}} \rfloor \leq N_{(2k+1,*)} \leq \lfloor 2^{k+1}\sqrt{2} \rfloor \end{aligned}$$

and

$$\begin{aligned} 2^{2k+3} < 2^{2k+3} + 1 \leq N_{(2k+2,*)} \leq 2^{2k+4} - 1 < 2^{2k+4} \\ \Rightarrow \lfloor 2^{k+1}\sqrt{2} \rfloor \leq \lfloor \sqrt{N_{(2k+2,*)}} \rfloor < 2^{k+2} \end{aligned}$$

Note that  $1 + \frac{1}{2} - \frac{1}{4} < \sqrt{2} < 1 + \frac{1}{2}$ , it yields

$$2^{k+1} + 2^k - 2^{k-1} \leq \lfloor 2^{k+1}\sqrt{2} \rfloor \leq \left\lfloor 2^{k+1}\left(1 + \frac{1}{2}\right) \right\rfloor = 2^{k+1} + 2^k$$

Hence it holds

$$2^{k+1} \leq N_{(2k+1,*)} \leq \lfloor 2^{k+1}\sqrt{2} \rfloor \leq 2^{k+1} + 2^k$$

and

$$2^{k+1} + 2^k - 2^{k-1} \leq \lfloor 2^{k+1}\sqrt{2} \rfloor \leq \lfloor \sqrt{N_{(2k+2,*)}} \rfloor < 2^{k+2}$$

That is  $(\lfloor \sqrt{N_{(2k+1,*)}} \rfloor \leq \lfloor 2^{k+1}\sqrt{2} \rfloor) \square k$  and  $(\lfloor 2^{k+1}\sqrt{2} \rfloor \leq \lfloor \sqrt{N_{(2k+2,*)}} \rfloor) \square k$ .

Considering the biggest node  $N_{(2k,2^{2k-1})}$  on level  $2k$ , it yields

$$N_{(2k,*)} \leq N_{(2k,2^{2k-1})} = 2^{2k+2} - 1 < 2^{2k+2} \Rightarrow \lfloor \sqrt{N_{(2k,*)}} \rfloor \leq \lfloor \sqrt{N_{(2k,2^{2k-1})}} \rfloor < 2^{k+1}$$

Likewise, considering the smallest node  $N_{(2k+3,0)}$  on level  $2k+3$ , it yields

$$N_{(2k+3,0)} = 2^{2k+4} + 1 > 2^{2k+4} \Rightarrow \lfloor \sqrt{N_{(2k+3,*)}} \rfloor \geq \lfloor \sqrt{N_{(2k+3,0)}} \rfloor \geq 2^{k+2} > 2^{k+2} - 1$$

Hence there is not a node  $N_{(2k,*)}$  satisfying  $N_{(2k,*)} \square k$ , and there is neither a node  $N_{(2k+3,*)}$  satisfying  $N_{(2k+3,*)} \square k$ .

**Proposition 3.** Node  $N_{(k,j)}$  ( $k > 1$ ) of  $T_3$  satisfies

$$2^{\lfloor \frac{k+1}{2} \rfloor} - 1 < 2^{\lfloor \frac{k+1}{2} \rfloor} \leq \lfloor \sqrt{N_{(k,j)}} \rfloor \leq 2^{\lfloor \frac{k+1}{2} \rfloor + 1} - 1 < 2^{\lfloor \frac{k+1}{2} \rfloor + 1}$$

Or equivalently  $\lfloor \sqrt{N_{(k,j)}} \rfloor \square \left\lfloor \frac{k-1}{2} \right\rfloor$ .

**Proof.** Since  $2^{k+1} + 1 \leq N_{(k,j)} \leq 2^{k+2} - 1$ , it yields  $2^{k+1} < N_{(k,j)} < 2^{k+2}$ ; hence it holds

$$2^{\frac{k+1}{2}} < \sqrt{N_{(k,j)}} < 2^{\frac{k}{2}+1}$$

By Lemma 2(P31), the inequality  $k \leq 2 \left\lfloor \frac{k+1}{2} \right\rfloor \leq k+1$  yields  $2^{\lfloor \frac{k+1}{2} \rfloor} \leq 2^{\frac{k+1}{2}}$  and  $2^{\frac{k}{2}+1} \leq 2^{\lfloor \frac{k+1}{2} \rfloor + 1}$ , hence it holds

$$2^{\lfloor \frac{k+1}{2} \rfloor} < \sqrt{N_{(k,j)}} < 2^{\lfloor \frac{k+1}{2} \rfloor + 1}$$

By Lemma 2 (P13) it immediately leads to

$$2^{\lfloor \frac{k+1}{2} \rfloor} \leq \lfloor \sqrt{N_{(k,j)}} \rfloor < 2^{\lfloor \frac{k+1}{2} \rfloor + 1}$$

which is  $\lfloor \sqrt{N_{(k,j)}} \rfloor \square \left( \left\lfloor \frac{k+1}{2} \right\rfloor - 1 = \left\lfloor \frac{k-1}{2} \right\rfloor \right)$ .

**Remark 2.** This proposition is a modification of the Corollary 2 in paper [3]. In paper [3], it claimed that  $\lfloor \sqrt{N_{(k,j)}} \rfloor \square \left\lfloor \frac{k+1}{2} \right\rfloor - 1$  or  $\lfloor \sqrt{N_{(k,j)}} \rfloor \square \left\lfloor \frac{k}{2} \right\rfloor$ . However, seeing from Propositions 1, 2 and 3, one can see that  $\lfloor \sqrt{N_{(k,j)}} \rfloor \square \left\lfloor \frac{k+1}{2} \right\rfloor$  never occurs.

**Theorem 2.** Given  $N > 3$  be an odd integer; then  $\lfloor \sqrt{N} \rfloor \square \left\lfloor \frac{\lfloor \log_2 N \rfloor}{2} \right\rfloor - 1$ .

**Proof.** Let  $k = \lfloor \log_2 N \rfloor - 1$ . By Lemma 1(P1),  $N$  is a node on level  $k$  of  $T_3$ . By Proposition 3 it knows  $\lfloor \sqrt{N} \rfloor \square \left\lfloor \frac{k-1}{2} \right\rfloor$ , that is  $\lfloor \sqrt{N} \rfloor \square \left\lfloor \frac{\lfloor \log_2 N \rfloor}{2} \right\rfloor - 1$ .

**Example 1.** Table 1 lists several odd integers that are randomly picked, and their positions in  $T_3$  as well as their square roots in  $T_3$ . It can see that,  $\lfloor \sqrt{N} \rfloor \square \left\lfloor \frac{\lfloor \log_2 N \rfloor}{2} \right\rfloor - 1$  holds for each number. Readers can check it manually or with Mathematica.

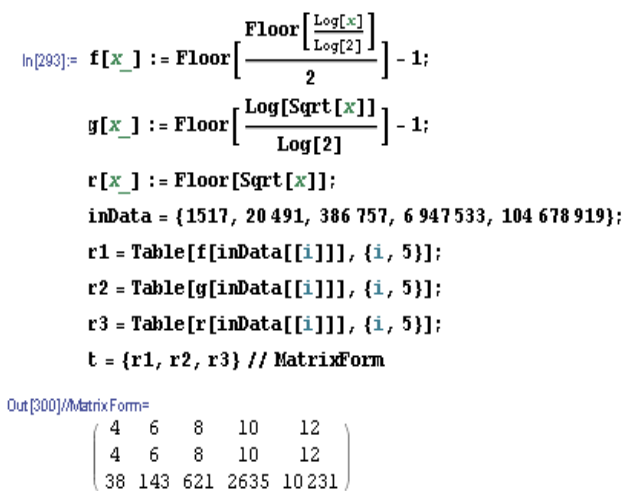
**Table1.** Odd Integers and its square roots in  $T_3$

Odd Integer $N$	$N$ 's in $T_3$	$\left\lfloor \frac{\lfloor \log_2 N \rfloor}{2} \right\rfloor - 1$	$\lfloor \sqrt{N} \rfloor$ & its level
1517	$N_{(9,246)}$	4	$\lfloor \sqrt{1517} \rfloor = 38 \square 4$
20491	$N_{(13,2053)}$	6	$\lfloor \sqrt{20491} \rfloor = 143 \square 6$
386757	$N_{(17,62306)}$	8	$\lfloor \sqrt{386757} \rfloor = 621 \square 8$
6947533	$N_{(21,1376614)}$	10	$\lfloor \sqrt{6947533} \rfloor = 2635 \square 10$
104678919	$N_{(25,18785027)}$	12	$\lfloor \sqrt{104678919} \rfloor = 10231 \square 12$

**Remark 3.** Table 1 can be easily checked manually or with Mathematica. When programmed as follows,

```
f[x_]:=Floor[Floor[Log[x]/Log[2]]/2]-1;
g[x_]:=Floor[Log[Sqrt[x]]/Log[2]]-1;
r[x_]:=Floor[Sqrt[x]];
inData={1517,20491,386757,6947533,104678919};
r1=Table[f[inData[[i]]],{i,5}];
r2=Table[g[inData[[i]]],{i,5}];
r3=Table[r[inData[[i]]],{i,5}];
t={r1,r2,r3} // MatrixForm
```

the screenshot in Mathematica 7.0 is as figure 1



**Fig1.** Screenshot of the program and output

**Corollary 1.** Let  $N > 3$  be an odd integer and  $k = \lfloor \log_2 N \rfloor - 1$ ; then  $\left(\lfloor \sqrt{N} \rfloor \leq \left\lfloor 2^{\lfloor \frac{k+1}{2} \rfloor} \sqrt{2} \right\rfloor\right) \square \left\lfloor \frac{k-1}{2} \right\rfloor$  when

$k$  is odd whereas  $\left(\left\lfloor 2^{\lfloor \frac{k+1}{2} \rfloor} \sqrt{2} \right\rfloor \leq \lfloor \sqrt{N} \rfloor\right) \square \left\lfloor \frac{k-1}{2} \right\rfloor$  when  $k$  is even.

**Proof.** By Theorem 2,  $\lfloor \sqrt{N} \rfloor \square \left\lfloor \frac{k-1}{2} \right\rfloor$  is sure. Now consider the fact that, whether  $k = 2l + 1$  or

$k = 2l + 2$ ,  $\left\lfloor \frac{k+1}{2} \right\rfloor = l + 1$  always holds. By Proposition 2, it knows that,  $\left(\lfloor \sqrt{N} \rfloor \leq \left\lfloor 2^{\lfloor \frac{k+1}{2} \rfloor} \sqrt{2} \right\rfloor\right) \square \left\lfloor \frac{k-1}{2} \right\rfloor$

when  $k$  is odd whereas  $\left(\left\lfloor 2^{\lfloor \frac{k+1}{2} \rfloor} \sqrt{2} \right\rfloor \leq \lfloor \sqrt{N} \rfloor\right) \square \left\lfloor \frac{k-1}{2} \right\rfloor$  when  $k$  is even.

#### 4. CONCLUSION

The square and the square root are important numbers. For an integer, the square root is essentially important because it is the cut-off point of divisors of a composite integer. Study of these numbers is helpful to understand distribution of the divisors on  $T_3$ . The Theorem 2 proved in this paper is of course a foundation for us to know where the square root of a node lies. Hope it is helpful in the future.

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**WANG Xingbo**, was born in Hubei, China. He got his Master and Doctor's degree at National University of Defense Technology of China and had been a staff in charge of researching and developing CAD/CAM/NC technologies in the university. Since 2010, he has been a professor in Foshan University, still in charge of researching and developing CAD/CAM/NC technologies. Wang has published 8 books, over 100 papers and obtained more than 30 patents in mechanical engineering

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